A Model of Type Theory with Directed Univalence in Bicubical Sets

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HoTT. August 14, 2019

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 - ∞-categories (Segal types) and univalent ∞-category (Rezk types) given internally as predicates on types
 - 4. Predicate isCov(B : A \rightarrow U) for covariant discrete fibrations
 - Cavallo, Riehl and Sattler have also (externally) defined the universe of covariant fibrations (the ∞-category of spaces and continuous functions) and shown *Directed Univalence:* Hom_{UCov} A B ≃ A → B

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 - 3. Use LOPS to define universe of covariant fibrations and construct directed univalence internally...
 - ...unfortunately, directed univalence is a bit trickier than expected

Let's see how far the techniques from cubical type theory get us!

Cubical Type Theory

(in the style of Orton-Pitts)

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I∷ Type

 $\mathbb{O}_{\mathbb{I}}$: \mathbb{I}

1∎:∎

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i.e. generators for the Cartesian cubes

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i.e. generators for the Cartesian cubes





Directed Type Theory

 $1_2:2$


Directed Type Theory

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<:2 y:2	x:2 y:2
x ∧ y : 2	x ∨ y : 2

and equations...

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- Why Dedekind cubes instead of Cartesian?
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- We also add the following axioms:
 - $p : \mathbb{I} \to \mathbb{Z}$ is constant ($\Pi x y : \mathbb{I}, p x = p y$)
 - $p : \mathbb{Z} \to \mathbb{Z}$ is monotone ($\Pi x y : \mathbb{Z}$, if $x \le y$ then $p x \le p y$)

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isCof : $\Omega \rightarrow \Omega$

Cof := $\Sigma \varphi$: Ω . isCof φ

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```
hasCom : (\mathbb{I} \to U) \to U
hasCom A = \Pi i j : \mathbb{I}.
\Pi a : Cof.
\Pi t : (\Pi x : \mathbb{I} . a \to A x)
\Pi b : (A i)[a \mapsto t i].
(A j)[a \mapsto t j; i = j \mapsto b]
```

```
relCom : (A : U) → (A → U) → U
relCom A B = Π p : \mathbb{I} \rightarrow A.
hasCom (B \circ p)
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Directed Type Theory

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```
hasCov : (2 \rightarrow U) \rightarrow U
hasCov A = \Pi \alpha : Cof .
\Pi t : (\Pi x : 2 \cdot \alpha \rightarrow A x)
\Pi b : (A \oplus_2)[\alpha \mapsto t \oplus_2] .
(A \mathbb{1}_2)[\alpha \mapsto t \mathbb{1}_2]
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• U_{Kan} given by LOPS construction for relCom

Directed Type Theory

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• U_{Kan} given by LOPS construction for relCom

Directed Type Theory

5. Define universe of covariant fibrations

 U_{Cov} given by LOPS construction for relCov.
 Lemma: relCov is in UKan

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• Key Idea: Glue type to attach equivalences to path structure

Directed Type Theory

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• Key Idea: Glue type to attach equivalences to path structure

Directed Type Theory

6. Construct directed univalence

 Key Idea: Glue type to attach functions to morphism structure





 $a \vdash Glue [a \mapsto (T, f)] B \equiv T$



 $\alpha \vdash Glue [\alpha \mapsto (T, f)] B = T$








Defining Directed Univalence

dua i A B f := Glue [i = 0 \mapsto (A , f : A \rightarrow B) , i = 1 \mapsto (B , id)] B : Hom_U A B



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 - dua_{β} : Π f : A \rightarrow B . Path f (dcoe (dua f))
 - dua_{nfun}: Πp : Hom A B. Πi : \mathbb{Z} . $p i \rightarrow$ (dua (dcoe p)) i

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- To complete directed univalence, we need $dua_{\eta fun}^{-1}$
- Agda: <u>https://github.com/dlicata335/cart-cube</u>

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- New goal: use any techniques available to confirm directed univalence holds at all in a cubical setting.
- Note: We would love any/all feedback on the math that follows.

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 - specifically: weak equivalences in the model are levelwise weak equivalences of simplicial sets

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 - specifically: weak equivalences in the model are levelwise weak equivalences of simplicial sets
- Dedekind cubes are not Reedy...

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 - tiny interval

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• The Idea: Categories permitting inductive constructions of presheaves and their morphisms (akin to cell complexes)

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- The Idea: Categories permitting inductive constructions of presheaves and their morphisms (akin to cell complexes)
- (informal/incomplete) Definition: A generalized Reedy category is a category C along with a degree function
 δ : ob C → N such that every morphism (that isn't an iso) factors through an object of strictly smaller degree

• Free Cartesian category on an interval generated by:

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 - face maps (+)

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 - diagonals (+)
- Free Cartesian category on an interval generated by:
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- Free Cartesian category on an interval generated by:
 - face maps (+)
 - diagonals (+)
 - degeneracies (-)
 - connections (-)

$(x, y) \mapsto (x, y, y) \mapsto (x \land y, y)$



$(\underline{x, y}) \mapsto (x, y, y) \mapsto (x \land y, y)$



$(x, y) \mapsto (\underline{x, y, y}) \mapsto (x \land y, y)$



up by a diagonal

$(x, y) \mapsto (x, y, y) \mapsto (\underline{x \land y, y})$



down by a connection











 $(x, y) \mapsto (x \land y, y)$



У

Х





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 - Inspired by Kapulkin and Voevodsky

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- Useful Lemma: We can build a topology J_{im} (the *image* covering) on Im(C) such that [C^{op}, Set] ≅ Sh(Im(C), J_{im}).
 - Inspired by Kapulkin and Voevodsky
 - The Comparison Lemma: [SGA 4, The Elephant]

• Definition: The *prism category* is the image closure of the Dedekind cube category.

- Definition: The *prism category* is the image closure of the Dedekind cube category.
- Lemma (with Christian Sattler): The prism category is equivalent to the full subcategory of simplicial sets containing subobjects of the Dedekind cubes (Γ, φ) generated by the following formulae:
 - ⊤ : true
 - x ≤ y : the equalizer of the degeneracy map x and connection x ∧ y
 - $\phi \land \psi$: the pullback of the subobjects (Γ , ϕ) and (Γ , ψ)
 - $\phi \lor \psi$: the pushout of the pullback for ($\Gamma, \phi \land \psi$)

• The Prism category

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 - is a finite product category...

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 - ...and thus the Yoneda embedding of its interval is tiny...

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 - The down maps are those that are regular epis in the presheaf category
 - The up maps are the monos
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- Corollary: The opposite of the prism category is also generalized Reedy

- Theorem: The prism category is a generalized Reedy category.
 - The down maps are those that are regular epis in the presheaf category
 - The up maps are the monos
 - The Reedy factorization is the image factorization
- Corollary: The opposite of the prism category is also generalized Reedy
- Question: For which categories C is Im(C) Reedy?

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- Can we make this even more cubical?
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- The lemma that finished directed univalence is still true after the transfer