# A Model of Type Theory with Directed Univalence in Bicubical Sets 

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HoTT. August 14, 2019

## Directed Type Theory

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- Riehl-Shulman defines a type theory for $\infty$-categories with a model bisimplicial sets


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4. Predicate isCov(B:A $\rightarrow \mathrm{U})$ for covariant discrete fibrations

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3. $\infty$-categories (Segal types) and univalent $\infty$-category (Rezk types) given internally as predicates on types
4. Predicate isCov(B:A $\rightarrow \mathrm{U}$ ) for covariant discrete fibrations
5. Cavallo, Riehl and Sattler have also (externally) defined the universe of covariant fibrations (the $\infty$-category of spaces and continuous functions) and shown Directed Univalence: Homucov A B $\sim A \rightarrow B$

## Constructive(?) Directed Type Theory

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- Can we make this constructive?


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1. Begin with Cubical Type Theory
2. Use a second cubical interval to define Hom-types
3. Use LOPS to define universe of covariant fibrations and construct directed univalence internally...

- ...unfortunately, directed univalence is a bit trickier than expected


# Let's see how far the techniques from cubical type theory get us! 

## Defining Bicubical Directed Type Theory

Cubical Type Theory (in the style of Orton-Pitts)

Directed Type Theory

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2. Add an interval: $\mathbb{I}$

## Directed Type Theory

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Cubical Type Theory (in the style of Orton-Pitts)
2. Add an interval: II

$$
\begin{array}{ll} 
& \mathbb{I}: \text { Type } \\
\mathbb{D}_{\mathbb{I}}: \mathbb{I} & \\
\mathbb{1}_{\mathbb{I}}: \mathbb{I}
\end{array}
$$

Directed Type Theory

2. Add an interval: $\mathbb{Z}$

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$\frac{\mathbb{I}: \text { Type }}{}$

$\frac{\mathbb{I}_{\mathbb{I}}: \mathbb{I}}{} \quad$| $\mathbb{1}_{\mathbb{I}}: \mathbb{I}$ |
| :--- |,

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2. Add an interval: $\mathbb{Z}$

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\overline{2}: \text { Type }
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Directed Type Theory

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\begin{gathered}
\overline{\mathbb{Z}: \text { Type }} \\
\frac{\mathbb{D}_{\mathbb{Z}}: \mathbb{Z}}{\mathbb{1}_{\mathbb{Z}}: \mathbb{Z}} \\
\frac{x: \mathbb{Z} \quad y: \mathbb{Z}}{x \wedge y: \mathbb{Z}}
\end{gathered} \frac{x: \mathbb{Z} y: \mathbb{Z}}{x \vee y: \mathbb{Z}}
$$

i.e. generators for the Cartesian cubes

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| :--- |,

Directed Type Theory

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\frac{x: \mathbb{Z} \quad y: \mathbb{Z}}{x \wedge y: \mathbb{Z}} \quad \frac{x: \mathbb{Z} \quad y: \mathbb{Z}}{x \vee y: \mathbb{Z}} \\
\text { and equations... }
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\frac{x: \mathbb{Z} \quad y: \mathbb{Z}}{x \wedge y: \mathbb{Z}} \quad \frac{x: \mathbb{Z} \quad y: \mathbb{Z}}{x \vee y: \mathbb{Z}} \\
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$$

i.e. generators for the Dedekind cubes

## The Directed Interval

- Why Dedekind cubes instead of Cartesian?

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x \leq y:=x=x \wedge y
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- $p: \mathbb{I} \rightarrow \mathbb{Z}$ is constant $(\Pi \times y: \mathbb{I}, p x=p y)$


## The Directed Interval

- Why Dedekind cubes instead of Cartesian?
$x \leq y:=x=x \wedge y$
- We also add the following axioms:
- $p: \mathbb{I} \rightarrow \mathbb{Z}$ is constant $(\Pi \times y: \mathbb{I}, p x=p y)$
- $p: \mathbb{Z} \rightarrow \mathbb{Z}$ is monotone $(\Pi x y: \mathbb{Z}$, if $x \leq y$ then $p x \leq p y)$


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Cubical Type Theory (in the style of Orton-Pitts)

## Directed Type Theory

1. Begin with Cubical Type Theory
2. Add an interval: $\mathbb{Z}$
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## Defining Bicubical Directed Type Theory

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3. Specify gen. cofibrations for II

$$
\begin{gathered}
\text { isCof }: \Omega \rightarrow \Omega \\
\text { Cof }:=\Sigma \phi: \Omega . \text { isCof } \phi
\end{gathered}
$$

Directed Type Theory

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Cof closed under _^_, _ی_, $\perp, ~ T$

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Cof closed under _^_, _ی_, $\perp, ~ T$

$$
\frac{x: \mathbb{I} \quad y: \mathbb{I}}{-: \text { isCof }(x=y)}
$$

$$
\begin{aligned}
& \phi: \mathbb{I} \rightarrow \text { Cof } \\
& . \text { in }
\end{aligned}
$$

$$
{ }_{-}: \text {isCof }(\Pi x: \mathbb{I} . \phi x)
$$

## Directed Type Theory

3. Specify gen. cofibrations for $\mathbb{Z}$

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$$
\begin{gathered}
\frac{x: \mathbb{Z} \quad y: \mathbb{Z}}{-: \text { isCof }(x=y)} \\
\phi: \mathbb{Z} \rightarrow \operatorname{Cof} \\
\hline: \text { isCof }(\Pi x: \mathbb{Z} \cdot \phi x)
\end{gathered}
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Directed Type Theory

4. Define filling problem for covariant fibrations

## Defining Bicubical Directed Type Theory

Cubical Type Theory (in the style of Orton-Pitts)
4. Define filling problem for Kan fibrations

$$
\begin{aligned}
& \text { hasCom: }(\mathbb{I} \rightarrow \mathrm{U}) \rightarrow \mathrm{U} \\
& \text { hasCom } A=\Pi i j: \mathbb{I} \text {. } \\
& \text { Пa: Cof. } \\
& \Pi t:(\Pi x: \mathbb{I} . a \rightarrow A x) \\
& \Pi b:(A i)[a \mapsto t i] . \\
& (A j)[a \mapsto t j ; i=j \mapsto b] \\
& \text { relCom: }(\mathrm{A}: \mathrm{U}) \rightarrow(\mathrm{A} \rightarrow \mathrm{U}) \rightarrow \mathrm{U} \\
& \text { relCom } A B=\Pi p: \mathbb{I} \rightarrow A \text {. } \\
& \text { hasCom ( } B \cdot p \text { ) }
\end{aligned}
$$

Directed Type Theory

4. Define filling problem for covariant fibrations

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4. Define filling problem for Kan fibrations

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\begin{aligned}
& \text { hasCom: }(\mathbb{I} \rightarrow \mathrm{U}) \rightarrow \mathrm{U} \\
& \begin{aligned}
\text { hasCom } A= & \Pi \mathrm{ij}: \mathbb{I} . \\
& \Pi \mathrm{a}: \operatorname{Cof} . \\
& \Pi t:(\Pi x: \mathbb{I} \cdot a \rightarrow A x) \\
& \Pi b:(A i)[a \mapsto t i] . \\
& (A j)[a \mapsto t j ; i=j \mapsto b]
\end{aligned}
\end{aligned}
$$

relCom: $(\mathrm{A}: \mathrm{U}) \rightarrow(\mathrm{A} \rightarrow \mathrm{U}) \rightarrow \mathrm{U}$
relCom $A B=\Pi p: \mathbb{I} \rightarrow A$. hasCom ( $B \cdot p$ )

## Directed Type Theory

4. Define filling problem for covariant fibrations

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\begin{aligned}
& \text { hasCov: }(\mathbb{Z} \rightarrow \mathrm{U}) \rightarrow \mathrm{U} \\
& \text { hasCov } A= \Pi a: \operatorname{Cof} . \\
& \Pi t:(\Pi x: \mathbb{Z} \cdot a \rightarrow A x) \\
& \Pi b:\left(A \mathbb{O}_{2}\right)\left[a \mapsto t \mathbb{D}_{2}\right] . \\
&\left(A \mathbb{1}_{2}\right)\left[a \mapsto t \mathbb{1}_{2}\right] \\
& \text { relCov: }(A: U) \rightarrow(A \rightarrow U) \rightarrow U \\
& \text { relCov } A B= \Pi p: \mathbb{Z} \rightarrow A . \\
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5. Define universe of Kan fibrations

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- Ukan given by LOPS construction for relCom


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5. Define universe of Kan fibrations

Directed Type Theory

5. Define universe of covariant fibrations

- Ukan given by LOPS construction for relCom
- Ucov given by LOPS construction for relCov. Lemma: relCov is in UKan


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6. Construct univalence

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Cubical Type Theory (in the style of Orton-Pitts)

Directed Type Theory

6. Construct directed univalence

- Key Idea: Glue type to attach equivalences to path structure


## Defining Bicubical Directed Type Theory

Cubical Type Theory (in the style of Orton-Pitts)
6. Construct univalence

- Key Idea: Glue type to attach equivalences to path structure

Directed Type Theory

6. Construct directed univalence

- Key Idea: Glue type to attach functions to morphism structure


## Glue Types

## Glue Types

Glue $[a \mapsto(T, f)] B:=$


## Glue Types



## Glue Types



## Glue Types




## Glue Types



$\frac{\mathrm{g}: \text { Glue }[\mathrm{a} \mapsto(\mathrm{T}, \mathrm{f})] \mathrm{B}}{\text { unglue } \mathrm{g}: \mathrm{B}}$ $a \vdash$ glue $t b \equiv t$

## Glue Types




## Glue Types



# Defining Directed Univalence 

dua i A B f:= Glue $[i=0 \mapsto(A, f: A \rightarrow B)$ $, i=1 \mapsto(B, i d)] B \quad: H o m \cup A B$


Naive Directed Univalence

## Naive Directed Univalence

- dua is Kan + covariant, and thus lands in UCov


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- dcoe : $(H o m A B) \rightarrow(A \rightarrow B)$


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- dcoe : $($ Hom A B) $\rightarrow(A \rightarrow B)$
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- duaß : П f: A $\rightarrow$ B. Path f (dcoe (dua f))


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- dua is Kan + covariant, and thus lands in $U_{\text {Cov }}$
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- dua: $(A \rightarrow B) \rightarrow$ Hom $A B$
- duaß : П f: A $\rightarrow$ B. Path f(dcoe (dua f))
- duanfun: П p:Hom A B . Пi: $\mathbb{Z} . p \mathrm{i} \rightarrow($ dua (dcoe $p)) \mathrm{i}$

Naive Directed Univalence

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- To complete directed univalence, we need duanfun ${ }^{-1}$


## Naive Directed Univalence

- We're thus left with the following picture:

- To complete directed univalence, we need duanfun ${ }^{-1}$
- Agda: https://github.com/dlicata335/cart-cube


## What next?

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- New goal: use any techniques available to confirm directed univalence holds at all in a cubical setting.


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- Cavallo, Riehl and Sattler's proof of directed univalence contains the precise lemma we need to finish.
- New goal: use any techniques available to confirm directed univalence holds at all in a cubical setting.
- Note: We would love any/all feedback on the math that follows.


## What next?

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- The proof in the bisimplicial model relies on simplices being a Reedy category


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- specifically: weak equivalences in the model are levelwise weak equivalences of simplicial sets
- Dedekind cubes are not Reedy...


## Our New Goal

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- Find a setting that...


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1. is cubical set valued presheaves of a Reedy category

## Our New Goal

- Find a setting that...

1. is cubical set valued presheaves of a Reedy category
2. interprets the axioms from our internal language

## Our New Goal

- Find a setting that...

1. is cubical set valued presheaves of a Reedy category
2. interprets the axioms from our internal language
3. allows for the LOPS construction of universes

## Our New Goal

- Find a setting that...

1. is cubical set valued presheaves of a Reedy category
2. interprets the axioms from our internal language
3. allows for the LOPS construction of universes

- tiny interval


## What are Reedy Categories?

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- The Idea: Categories permitting inductive constructions of presheaves and their morphisms (akin to cell complexes)
- (informal/incomplete) Definition: A generalized Reedy category is a category C along with a degree function $\delta:$ ob $\mathrm{C} \rightarrow \mathbb{N}$ such that every morphism (that isn't an iso) factors through an object of strictly smaller degree


## The Dedekind Cubes

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- Free Cartesian category on an interval generated by:


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- Free Cartesian category on an interval generated by:
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- Free Cartesian category on an interval generated by:
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up by a diagonal

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down by a connection

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- Useful Lemma: We can build a topology Jim (the image covering) on $\operatorname{Im}(C)$ such that $[\mathrm{Cop}, \mathrm{Set}] \cong \mathrm{Sh}(\operatorname{lm}(C), \mathrm{Jim})$.


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- The Comparison Lemma: [SGA 4, The Elephant]


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- Lemma (with Christian Sattler): The prism category is equivalent to the full subcategory of simplicial sets containing subobjects of the Dedekind cubes ( $Г, \Phi)$ generated by the following formulae:
- T: true
- $x \leq y$ : the equalizer of the degeneracy map $x$ and connection $x \wedge y$
- $\phi \wedge \psi$ : the pullback of the subobjects $(\Gamma, \Phi)$ and $(\Gamma, \psi)$
- $\phi \vee \psi$ : the pushout of the pullback for $(\Gamma, \phi \wedge \psi)$


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- ...and thus the Yoneda embedding of its interval is tiny...


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- Corollary: The opposite of the prism category is also generalized Reedy
- Question: For which categories $C$ is $\operatorname{Im}(\mathrm{C})$ Reedy?


# Model Category One: Prismatic Cubical Sets 

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- As our internal language axioms interpret into this model, we get a model with directed univalence!
- Can we make this even more cubical?


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- Left Induced Model Structure: [Hess-Kedziorek-RiehlShipley, Garner-Kedziorek-Riehl]
- Path Object Argument: [Quillen]
- Our internal language axioms still interpret after the transfer
- The lemma that finished directed univalence is still true after the transfer

